

Abstract

This presentation addresses the approaches, values, thresholds, triggers, and relationships of performance data analysis. Performance monitoring, analysis, and tuning should be a pro-active and continuous process; however, this often gets neglected because people are too busy with other work, or they are not sure what data is really important and which data quickly shows that there are problems - when nobody is complaining. Some basic thresholds and data provide quick insights about problems both at the application level, and the system level. Charts and graphs illustrate the surprising CPU and dollar costs of sequential prefetch and IOs.

This presentation does <u>not</u> address SQL coding or application design considerations.

Major Bullet Points

- Sources of data
- Key performance thresholds and indicators
 - Do you have problems?
- Rules of Thumb
 - Good numbers, bad numbers, or ???
 - Can't dot every *i* or cross every *t* in this presentation....
- Using and applying data for analysis
- Performance examples from the real world

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Outline:

- 1. DB2 Data Sources
 - a. Application Accounting Data
 - b. System Statistics Data
 - c. DB2 performance trace data
- 2. The important performance variables and indicators:
 - a. Variables and relationships
 - b. Inter-relationships between Statistics and Application data
 - c. How do you know when you have a problem, and just how bad is it?
- 3. Top down tuning approach
- Tuning from available DB2 data sources
 - b. How much can you tune the payback potential
 - c. How long will it take?
- 5. Bottom up approach
 - a. Data interpretation
 - b. Determining where the benefits are...
- 6. Summary, Guidelines and recommendations

We care about performance because..

- Competitiveness in the marketplace
- Providing service to your customers
 - · Both internal & external
- Reducing costs/avoiding costs
 - Tune, or buy hardware
 - Tuning is a free payback for the money already invested
- Time is money
 - The system is slow today... I can fix that...

Or maybe you're just a computer geek like me...

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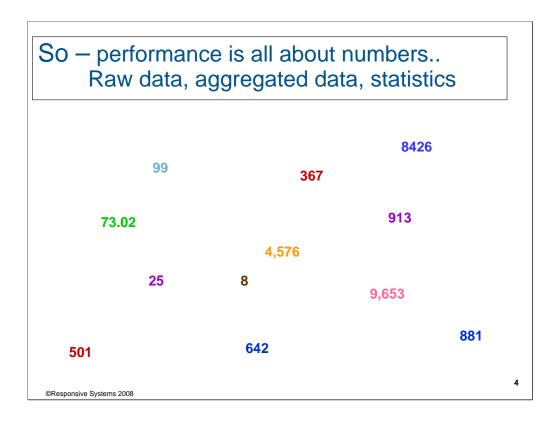
In today's competitive economic environment we can't be complacent about performance. Information is critical to the success of your company, and the ability to both retain existing customers and get new ones, may determine long term growth. The enterprises that can deliver information quickly will grow and thrive - those that cannot, may either cease to exist or will be swallowed by other more aggressive corporations.

Ever increasing processor capacities are often driven more by application inefficiencies and other performance related problems than by true volume demands. Throwing multi-million dollar hardware solutions at performance problems has become the norm.

If corporate stockholders became aware of the magnitude of corporate waste and inefficiencies, many corporate officers might find themselves on the receiving end of legal actions - and many boards of directors would be replaced.

Some reasonable amounts of system and application tuning can provide dramatic paybacks at many companies. The multi-million dollar tuning success stories exist and should receive more press than the large system failures.

If management could only become truly aware of the opportunities, and reward staff for improving performance - they would be amazed at the results and long term savings.



Numbers, numbers everywhere. What do they all mean? What is important?

Mean and Median

• How many of you remember the difference?

Mean is an average, and averages are easy - right?...

Mean is an average, and averages are easy – misleading...

$$4, 8, 12, 2, 10, 8, 104, 4, 6, 10 = 168/10 = 16.8$$

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Mean and Median

Median is a mid-point

2, 27, 11, 19, 33, 22, 4, 2, 19, 24, 22, 40, 38

2, 2, 4, 11, 19, 19, **22**, 22, 24, 27, 33, 38, 40 odd # values

2, 2, 4, 11, 19, 19, **22, 24**, 27, 33, 38, 40

even # values

23 22+24=46/2

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What does this mean to us?

- · We need multiple data or focus points
 - · Drill down to lower levels of detail
 - Exception reporting should be pushed from detail upwards....
- Long averages across periods mask/hide problems
 - - Average is 12.4 but it's 24% higher than the median, or the mode

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DB2 Performance is relative to work

- Work performed by the application
 - Getpages
 - Fetches
 - I/Os
- · Class 1 elapsed time
 - · Ratio compared to class 2 elapsed time
- Class 2 elapsed time (online transactions)
 - .050 seconds
 - .100 seconds
 - .500 seconds

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Performance is relative. What are your expectations? Are they reasonable? What constitutes good performance? What is a reasonable amount of work for an online transaction? How do we characterize this?

Many well designed online transaction applications have average elapsed times of less than .020 seconds, and provide true sub-second response time to the end user. However, applications like this are becoming exceedingly rare in today's environment of bloated application code, poorly designed databases, and poorly coded SQL.

We will try to address what is good, what is reasonable, and basic metrics that indicate when performance is poor or marginal.

Keys to better performance

- Thought
- Inquisitiveness
- Reasonable
- Expectations
- Monitoring
 - Test/Dev Early Warning
- Performance History
 - Yesterday?
 - Last week?

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- Data relationships
- Knowing what is good, or bad
 - And why
- Caring about it
- Not waiting for complaints
- It can Always be Better

Performance is everyone's job, and responsibility!

Simply caring about performance issues, and looking at data will provide opportunities to make things better.

Those who wait for complaints, or wait for someone else to take action are not serving their company well - they are having a negative impact on the profitability of the entire corporation.

When you find something, and fix it - tell the world!!

Tell them how much you saved them, and how many more opportunities exist in your systems and applications.

Delaying or avoiding a processor upgrade is worth millions!!

Top down tuning

- Important/Meaningful periods
- Summary data don't get buried in detail
- What looks big, or strange, or not like most other things
- Don't assume others are doing their job properly
- Use multiple tools to look at performance

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It is not usually helpful to look at reports summarized over an entire days, or several days. You should be primarily concerned with the peak workload periods - if you can process well there, the rest of the time will run just fine....

Don't assume that others are doing their jobs, or even that they always know what they are doing - unless you are confident of a persons ability. Even then, they may be overloaded and not have the time do everything you think they are doing. Priorities of work assigned to others changes at the discretion of their management, so unless you work for the same person you can't always be sure what the priorities are.

Different tools often use a slightly different perspective to look at the same set of data - so seek input from all tools available to you.

Top down tuning

- Know what values are good, what is OK, what is bad
- Know what is normal for your system
- Exception reporting
- Don't get too bogged down in detail...
- But be able to drill down to find the lower level data

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It is really important to understand what performance numbers are good, and then what is <u>normal</u> for your system and applications. Tracking performance over time is a great way to find problems before they become critical.

Performance is like a traffic circle

- Application problems may lead to the DB2 system
- DB2 system problems often lead to application design and SQL coding
- Application and system problems may lead to DASD, z/OS, CICS, IMS
- z/OS performance problems may lead to applications or DB2 memory consumption



Memory and paging V8 and V9..

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As you find a problem, as correct it, other performance problems will surface that may not have been obvious before.

Monitoring and tuning is a constant process, not a one time exercise.

The only exception to this is if/when your environment is static - no application changes, no workload or data growth, no migration to newer releases or maintenance levels....

System Overhead

- (C2 Elapsed C3 Wait) / C2 CPU
- So what's a reasonable ratio?
 - Online IMS/CICS transactions

Batch

10.0 + or -

< 1.5

- These are affected by:
 - · Processor busy rate, possible system paging
 - Address space priorities, WLM setup

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An area that is often overlooked is the overall system overhead. If the priorities for all the address spaces do not have the correct relationship, varying amounts of overhead will slow your work.

Is your processor more than 95% busy on a regular basis? If so, your performance is hurt by the machine busy rate. Running at 100% is not a case of using the machine to its fullest capacity - it's a case of hurting all performance because the processor is overloaded!

Important application performance data

- Workload
- Class 1 2 relationship
- Class 3 Wait time
 - Components
 - Avg I/O times
 - Total % of Class 2
- Deadlocks
- Timeouts

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- Buffer pool performance
 - I/O per second (important)
 - Type of I/Os
 - Application hit % (useless)
 - System hit % (historical (GP-PagesRd)/GP & not useful)
- Ridpool failures
 - How many
 - Why
- Online transaction Seq. scan

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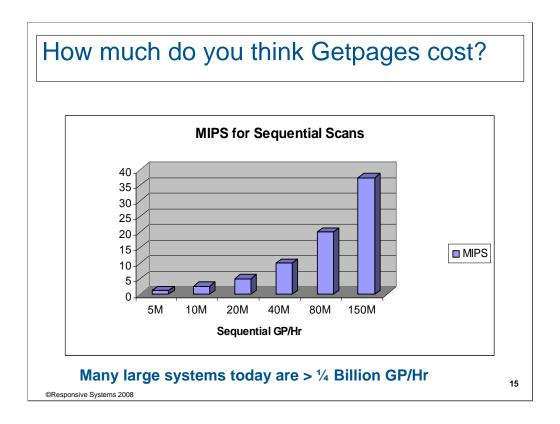
You must always consider the amount of work your application is performing - look at the number of getpages, fetches, I/Os, locks, etc, etc.

Look for the problem areas. What can be improved?

Calculate the System BP IO Rates, calculate the avg. I/O wait time by dividing the total synch wait time by the number of I/Os.

Look at any failures of resources, and look for things that don't seem normal to you.

Much of this can be programmed into your online monitor exception reports, or other reporting facilities. You really want a reasonable level of exception reporting so you don't have to look at tons of detail reports.



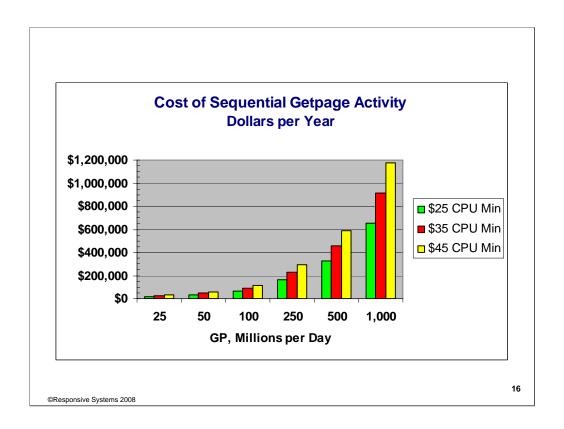
Finding one object dominating your system with sequential I/O?

One company fixed a problem like this years ago, and got back 18% of the machine and cancelled their upgrade!

Think you don't have sequentially accessed data in your system? Don't have some indexes with sequential scan?

Think again.... Most installations do – and nobody is complaining about performance. But management is complaining about CPU consumption and having to upgrade.....

But they won't spend a nickel to tune!!!

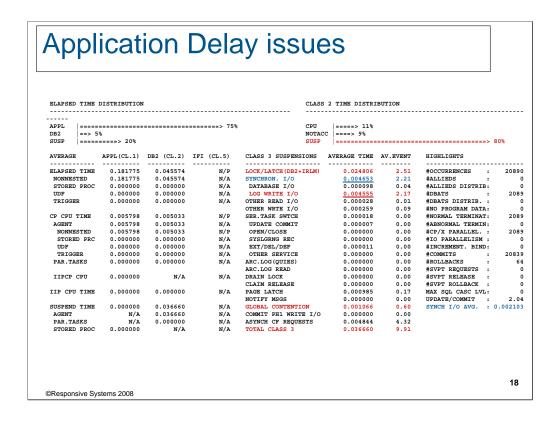


This is based on a 2064 processor with 210 MIP engine speeds.

Appl	icat	ion	Dela	ay iss	ue	S				
LOCATION: CASI GROUP: DBDG MEMBER:		C		R DB2 PERFORMANCE NG REPORT - LONG	EXPERT	(V3)	REQUESTED	PAGE: 1-4) FROM: 06/14/07 20:0 TO: 06/14/07 20:0		
SUBSYSTEM:			ORDER: PLANNAME-REQLOC				INTERVAL FROM: 06/14/07 20:05:25.65			
DB2 VERSION:	V8		SCOPE:	MEMBER				TO: 06/14/07 20:08:5	9.11	
PLANNAME: DIS ELAPSED TIME	DISTRIBUTION	I			CLASS 2	TIME DISTRIE	BUTION			
	*				CPU NOTACC SUSP	======> =====> 12%		> 67%	**:	
AVERAGE		DB2 (CL.2)	TET (CT. 5)	CLASS 3 SUSPENS		AVERAGE TIME		HIGHLIGHTS		
ELAPSED TIME	0.139919	0.018450	N/P	LOCK/LATCH (DB2+1	IRLM)	0.003620	1.28	#OCCURRENCES :	33'	
NONNESTED		0.018450		SYNCHRON. I/O		0.005955	2.12	#ALLIEDS :		
STORED PROC		0.000000		DATABASE I/O		0.000017	0.02	#ALLIEDS DISTRIB:		
UDF	0.000000	0.000000		LOG WRITE I/O		0.005938	2.11	#DBATS :	33	
TRIGGER	0.000000	0.000000	N/A	OTHER READ I/O		0.000006	0.00	#DBATS DISTRIB. :		
			N/P	OTHER WRTE I/O		0.000321	0.18	#NO PROGRAM DATA:		
CP CPU TIME AGENT	0.004498	0.003869		SER.TASK SWTCH UPDATE COMMIT		0.000016	0.00	#NORMAL TERMINAT: #ABNORMAL TERMIN:	33	
NONNESTED	0.004498	0.003869		OPEN/CLOSE		0.000002	0.00	#CP/X PARALLEL. :	33	
STORED PRC		0.000000		SYSLGRNG REC		0.000000	0.00	#IO PARALLELISM :	3.	
UDF	0.000000	0.000000		EXT/DEL/DEF		0.000014	0.00	#INCREMENT. BIND:		
TRIGGER	0.000000	0.000000	N/A	OTHER SERVICE		0.000000	0.00	#COMMITS :	336	
PAR.TASKS	0.000000	0.000000	N/A	ARC.LOG(QUIES)		0.000000	0.00	#ROLLBACKS :		
				ARC.LOG READ		0.000000	0.00	#SVPT REQUESTS :		
IIPCP CPU	0.000000	N/A	N/A	DRAIN LOCK		0.000000	0.00	#SVPT RELEASE :		
				CLAIM RELEASE		0.00000	0.00	#SVPT ROLLBACK :		
IIP CPU TIME	0.000000	0.000000	N/A	PAGE LATCH		0.000151	0.06	MAX SQL CASC LVL:		
				NOTIFY MSGS		0.000000	0.00	UPDATE/COMMIT :	1.	
SUSPEND TIME	0.000000	0.012340		GLOBAL CONTENTIO		0.000185	0.22			
AGENT	N/A		N/A	COMMIT PH1 WRITE			0.00			
PAR.TASKS STORED PROC	N/A		N/A	ASYNCH CF REQUES	STS	0.002085	2.94			
STORED PROC	0.000000	N/A N/A	N/A N/A	TOTAL CLASS 3		0.012340	6.80			
ODE	0.00000	N/A	N/A						1	

Within this five minute interval, there were 33,639 commits, and 64 ROLLBACKS.

Note that 67% of the application elapsed time is $\it WAIT$. Log Write IO is 48% of the Class 3 Wait Time!



Hopefully obvious the big problem is locking....

Little Red Flags...

SUBSYSTEM SERVICE COMPONENT IDENTIFY 31529 CREATE THREAD 757527 SIGNON 28235 TERMINATE 789499 ABORT 180102 COMMIT PHASE 1 3546312 COMMIT PHASE 2 2910873 READ ONLY COMMITS 643648 UNITS OF REC GONE INDOUBT 0 UNITS OF REC INDOUBT RESOLVED 0 SYNCHS (SINGLE PHASE COMMIT) 1857513

Hopefully, it should be obvious that large numbers of aborts can be a problem.

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```
Little Red Flags...
   From 03/28/08 09:00:00 To 03/28/08 09:30:00
  EXCEPTIONS Crit Warn Info

        Subsystem
        11
        2
        16
        EDET FOR Pree Pg
        20687
        % Total
        41.4

        ADDRESS SPACE CPU
        Warnings
        2440
        DBD Lds
        0
        % Rqsts
        0.0

        DBAS 10:24.95
        Act Pools
        0
        CT Lds
        0
        % Rqsts
        0.0

        SSAS 55.92
        % NStl Pgs
        7.8
        PT Lds
        26
        % Rqsts
        0.0

        IRLM 2.70
        Getpages 135233941
        Dyn Ins
        0
        % Rqsts
        0.0

        DIST 55:00.61
        Sync Rds
        5296912
        LOCKING
        LOGGING

        POOL FAILURES
        Buf Updts
        8210578
        Suspend
        2629
        Dlyd Wrts
        0

        RID
        0
        Pg Writes
        937334
        Escalate
        0
        Arch Read
        0

        EDM
        0
        Write I/O
        76331
        Timeout
        0
        Min/Ckpt *****

        THREADS
        SQL
        DATA SHARING
        STORED PROCS

                                                                                                                                                          EDM POOL
      Subsystem 11 2 16
                                    44345
                                                                                                                              25710
      Created
                                                                                    Dynamic
                                                                                                                                                                  Group DSNDBXX CALLs
                                                                                                                                                                                                                                                       352
                                                                                     In+Up+Dl 2171956
     Terminated 44653
                                                                                                                                                                  Member DB66
                                                                                                                                                                                                                             Fails
     Aborts
Commits
                                                    5986
                                                                                      Open+Sel 1153851
                                                   84603
                                                                                                                                                                                                                                                                               20
```

Hopefully, it should be obvious that large numbers of aborts can be a problem. At this time, there's a lot of free space in the EDM pool.

This should be tracked across other time frames to see if memory is wasted.

DB2 Read I/Os

Two basic types of object access and resulting I/Os

- Random or Synchronous I/O
- Prefetch or Asynchronous I/O
 - Sequential
 - Dynamic
 - List Prefetch

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There are many thresholds and exceptions that should be monitored and tracked on a regular basis, and some are more critical than others.

SPTH, DMTH, WITH, and EDMPool failures are critical and require immediate action - IWTH is not a problem unless you have hit SPTH and DMTH first.

Log Buffer unavailable and Unavailable Read/Write engines often points to poor dasd subsystem performance, and may have to be addressed from both sides to eliminate the problem.

Buffer Pool – Tuning Methodology

- Mostly Random RAMOS
 - · Small/Medium Working sets
 - Large Working sets
 - Very Large & Very Random
 - Never have a good Hit %
 - May have high I/O rate
 - Indexes
- Mostly Sequential SAMOS
 - Small/Medium Working sets
 - May be able to get a decent Hit % Large Working sets
 - Never get a decent Hit %
 - Don't need a large pool

WkSet is subjective – every system and application is different

Change the size of a pool, and every object wkset changes almost....

Move an object to a different pool, and the WkSet size changes

You must be able to <u>predict the</u> <u>effect of changes</u>, or you are guessing with your performance life

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A proven pool tuning methodology is the proper grouping of objects based on access type and working set size (number of resident pages in the pool). Object catalog statistics indicating the number of physical pages are not useful for this approach.

An object may have a million physical pages, but the important thing is how many you re-reference within a few minutes – and the impact this object reference pattern has on other objects in the pool and vice versa.

Buffer Pool – Tuning Methodology

- What is a transient object?
 - · Large wkset size
 - Wkset grows proportionally to pool size increases
 - · Pages not frequently re-referenced
- When do we care?
 - · When it impacts other objects in the pool
 - If it has a low GP and IO rate, we don't care
 - · High, we care a lot and immediately..
- Typically very large, and random or sequential
 - Random may have low or high Dynamic Prefetch...

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We don't have any control over dynamic prefetch. This is determined by the buffer manager, at the application cursor level. It may help performance of the object in use, and it may also hurt other objects in the pool.

Pool Performance

- Is rarely linear
 - · More memory does not always improve performance
 - Small pool increases often don't show any improvement
 - Sometimes additional "smallish" increases can provide substantial improvements
 - Many times large increases do not help
 - Sometimes they do....
 - It may depend on what "large" means to you...
 - A number, percentage, or it depends...

DB2 Performance – it depends...

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We will see from data later in this presentation that pool performance is not linear. Doubling the pool size does not double the hit ratio, or cut the I/O rate in half. We will always reach a point of diminishing returns, when adding buffers to a pool.

Now, I realize that some of the above items seem contradictory, and will explain them in more detail during the presentation.

The data shown in future slides will also illustrate all the above points.

Performance tuning complaint...

"I modified a bufferpool to set the DWQT from the default 50% value to DWQT=4 and VDWQT=0 - No other change -

After the change, a program doing the inserts (in ascending key), and delete on it, took a lot of time, 40min instead of 17min (it is a temporary table), STROBE shows that 90% of the time was on the Insert, and 80% of the Wait was on "OTHER WRITE".

I proved to them that the INDEX and TS has never been organized, and an increase in volume can magnify the problem. (Statistics showed 70,000 inserts instead of 10,000 during this day)

That's 7 times the workload, but only 2.3 times the elapsed

They told me that the problem comes from the change, because even if it was disorganized, other executions were good ... "

Three blind mice....

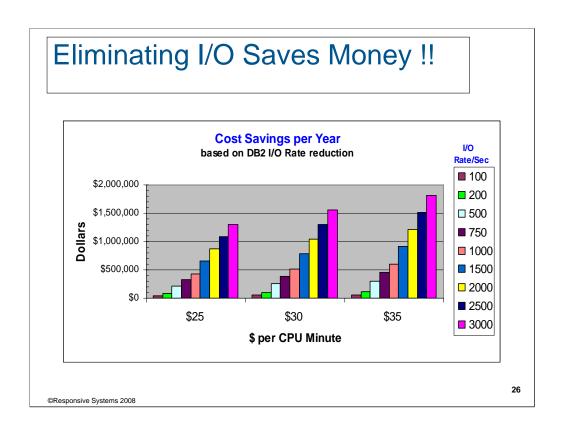
History data, history data...

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Wow, do you think that increasing the workload by a factor of 7 has any impact? Actually, performance was much better, not worse, and not longer.

The elapsed time was only 2.3 times the original job elapsed times.



These I/O rate per sec savings, up through 2,500 per second, have been achieved by clients

Buffer Pool Performance Data - Red Flags

BP5 GENERAL				/COMMIT
CURRENT ACTIVE BUFFERS UNAVAIL.BUFFER-VPOOL FULL	410.10	N/A	N/A	
NUMBER OF DATASET OPENS	8446.00	0.10	0.01	0.00
BUFFERS ALLOCATED - VPOOL	7507.00	N/A	N/A	N/A
DFHSM MIGRATED DATASET	0.00	0.00	0.00	0.00
DFHSM RECALL TIMEOUTS	0.00	0.00	0.00	0.00
VPOOL EXPANS. OR CONTRACT.	0.00	0.00	0.00	0.00
VPOOL OR HPOOL EXP.FAILURE	0.00	0.00	0.00	0.00
CONCUR.PREF.I/O STREAMS-HWM	300.00	N/A	N/A	N/A
PREF.I/O STREAMS REDUCTION	3382.00	0.04	0.00	0.00
PARALLEL QUERY REQUESTS	88003.00	1.02	0.12	0.03
PARALL.QUERY REQ.REDUCTION	215.00	0.00	0.00	0.00
PREF.QUANT.REDUCED TO 1/2	866.2K	10.06	1.21	0.34
PREF.QUANT.REDUCED TO 1/4	73224.00	0.85	0.10	0.03

What critical data items are missing from these sets of data?

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There are a lot of red flags in this data report, all related to a lack of buffers available for prefetch.

Buffer Pool Pe	311011	nanc	e Da	สเล -	Red Flags
BP5 READ OPERATIONS	QUANTITY	/SECOND	/THREAD	/COMMIT	
BPOOL HIT RATIO (%)	31.59				
GETPAGE REQUEST		2083.30	250.37	70.40	
GETPAGE REQUEST-SEQUENTIAL					4
	49309.8K				
SYNCHRONOUS READS	5285.0K	61.38	7.38	2.07	
SYNCHRON. READS-SEQUENTIAL	1797.4K	20.88	2.51	0.71	
SYNCHRON. READS-RANDOM	3487.7K	40.51	4.87	1.37	
GETPAGE PER SYN.READ-RANDOM	14.14				
SEQUENTIAL PREFETCH REQUEST	3856.3K	44.79	5.38	1.51	
SEQUENTIAL PREFETCH READS	3684.0K	42.79	5.14	1.45	
PAGES READ VIA SEQ.PREFETCH	109.6M	1273.45	153.05	43.04	
S.PRF.PAGES READ/S.PRF.READ	29.76				
LIST PREFETCH REQUESTS					
	392.2K			0.15	
PAGES READ VIA LIST PREFTCH	2049.5K	23.80	2.86	0.80	
L.PRF.PAGES READ/L.PRF.READ	5.23				
DYNAMIC PREFETCH REQUESTED	197.4K	2.29	0.28	0.08	
DYNAMIC PREFETCH READS	184.4K	2.14	0.26	0.07	
PAGES READ VIA DYN.PREFETCH	5726.7K	66.51	7.99	2.25	
D.PRF.PAGES READ/D.PRF.READ	31.06				
PREF.DISABLED-NO BUFFER	0.00	0.00	0.00	0.00	
PREF.DISABLED-NO READ ENG	74.00	0.00	0.00	0.00	—
PAGE-INS REQUIRED FOR READ	0 00	0.00	0.00	0.00	

Running out of Read Engines is often a sign of poor dasd performance.

Buffer Pool Performance Data - Red Flags

BP5 WRITE OPERATIONS	QUANTITY	/SECOND	/THREAD	/COMMIT
	10460 5	014 41	05.55	
BUFFER UPDATES	18460.7K	214.41	25.77	9.39
PAGES WRITTEN	2220.6K	25.79	3.10	1.13
BUFF.UPDATES/PAGES WRITTEN	8.31			
SYNCHRONOUS WRITES	25517.00	0.30	0.04	0.01
ASYNCHRONOUS WRITES	1063.6K	12.35	1.48	0.54
PAGES WRITTEN PER WRITE I/O	2.04			
HORIZ.DEF.WRITE THRESHOLD	1.00	0.00	0.00	0.00
VERTI.DEF.WRITE THRESHOLD	451.00	0.01	0.00	0.00
DM THRESHOLD	0.00	0.00	0.00	0.00
WRITE ENGINE NOT AVAILABLE	85.00	0.00	0.00	0.00
PAGE-INS REQUIRED FOR WRITE	0.00	0.00	0.00	0.00

Running out of read engines & write engines is often a sign of DASD performance problems. 600 read engines, 300 write engines (apar).

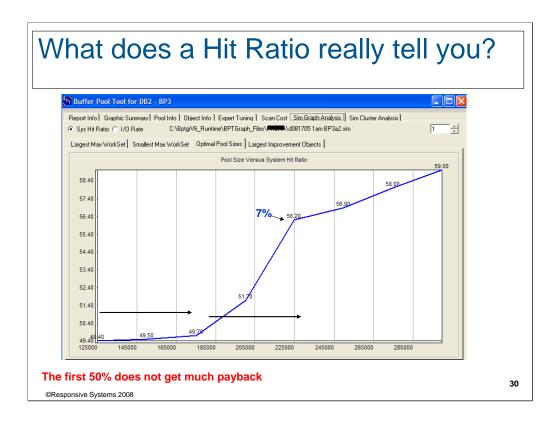
What critical data items are missing from these sets of data?

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The number of buffers in the pool, and pool thresholds.. And the ELAPSED TIME for the data!!

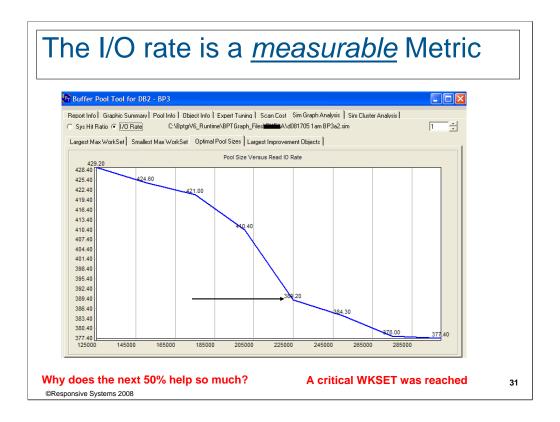
V8 has 600 read and 300 write engines; however, this number is for the entire DB2 system, not one pool.



Ok, it shows you that performance is better. But how much better is it? How much CPU and elapsed times have been saved from I/O avoidance?

Increasing the pool by 50% does not give much payback, the next 50,000 shows a large improvement, and then the improvement curve flattens.

Again, it looks nice, but you can't take any of the numbers to the bank.



The I/O rate is convertible into CPU costs, and elapsed time savings.

This is not just a suggestion to make the pool larger, it shows you the real benefit, and where to stop.

It shows you that the first 50,000 additional buffers don't provide much payback, but the next 50,000 give a huge payback.

The large payback from the second increment of 50,000 buffers is because we passed a critical working set threshold for a heavily accessed object. As stated earlier, the wkset size of an object has nothing to do with the number of pages shown in the catalog. It is the number of pages in the pool at a specific point in time.

The I/O rate is a *meaningful* Metric

INTV	INTV	GET	SYNC	SEQ	LIST	DYN	IORATE	HIT
DATE	TIME	PAGE	IO	PREFETCH	PREFETCH	PREFETCH	/SEC	RATIO
2007-02-2	1 00.00.00	1,187,068	43,331	625,280	156,611	202	459	30.47% ++++
2007-02-2	1 00.30.00	10,913,350	342,348	3,226,519	289,002	1,069,556	2,737	54.85%
2007-02-2	1 01.00.00	671,955	51,854	373,032	49,730	893	264	29.24%
2007-02-2	1 01.30.00	743,700	67,202	435,841	2,948	95	281	31.95%
2007-02-2	1 09.30.00	4,105,758	124,423	2,961,328	12,364	69,313	1,760	22.85%
2007-02-2	1 10.00.00	4,232,240	100,715	3,088,771	9,277	64,952	1,813	22.88%
2007-02-2	1 10.30.00	3,127,959	70,646	2,452,744	18,310	58,303	1,444	16.88% ++++
2007-02-2	1 11.00.00	3,890,741	112,328	2,520,929	9,911	51,934	1,497	30.73%
2007-02-2	1 11.30.00	4,011,848	117,173	2,394,283	11,760	74,255	1,443	35.26%
2007-02-2	1 12.00.00	4,580,930	109,663	2,339,369	639,288	68,607	1,754	31.09%
2007-02-2	1 12.30.00	6,150,020	188,803	3,202,333	11,852	64,002	1,926	43.63%

If the hit ratio was *meaningful*, it would not show a big increase when there is a large increase to the IO rate % GP increase vs. IO yes we did find a lot more of the pages in the pool....

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The IO rate can increase, and the hit ratio can increase.

The IO rate can decrease, and the hit ratio can decrease.

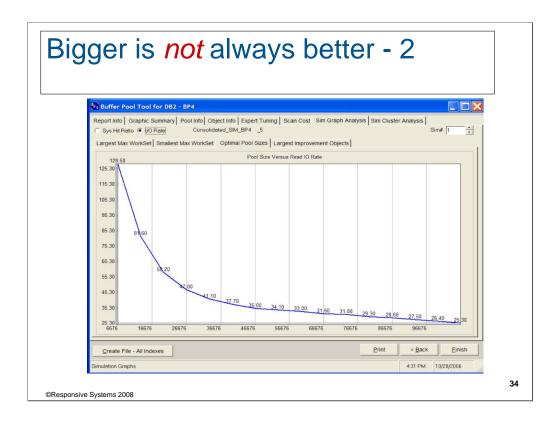
This is the opposite the "expectation"

Changes in the workload, the type of accesses taking place, and the objects in use, cause the counter-intuitive swings of the hit ratio. But it's counter-intuitive only if you are looking uniquely at getpages and IOs.



The increase for the buffer pool hit ratio flattens, and drops to .4% for every 10,000 buffers, 40 meg of memory.

As stated many times, these gains cannot be equated to elapsed times or CPU reductions.



The IO reduction/gain from increasing the pool size flattens, and eventually drops to only 1 IO/Sec per 10,000 buffers. This is not considered a useful payback for 40 Meg of memory.

18,000 to 36,000 saves about 15 IO/Sec., and we can see how the rate of saving drops off to almost nothing as we continue to add more memory.

Overall, the 36-38,000 buffer range is the best range for performance/memory trade-off.

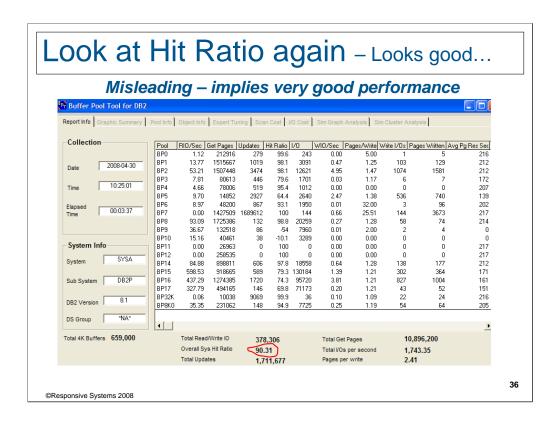
The previous slide with graph showed...

- When a pool is much too small, more memory will provide substantial improvements
- There is a point of limited, and possibly no return
 - Further increases provide very little gain, not enough to justify the added memory
 - The first tripling of size cut the IO rate 50%, 70 IOs
 - Then a doubling cut the IO rate by 15 IOs
 - The third doubling cut the IO rate by 5 IOs

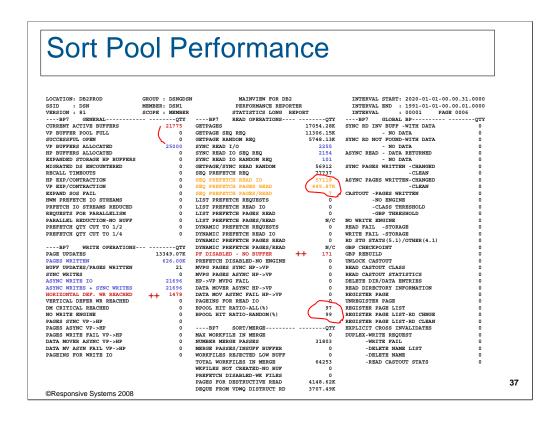
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Roughly, increase 6,000 to 18,000 has a large payback -70 IO/Sec, and is a 54% saving of the IO rate/sec.



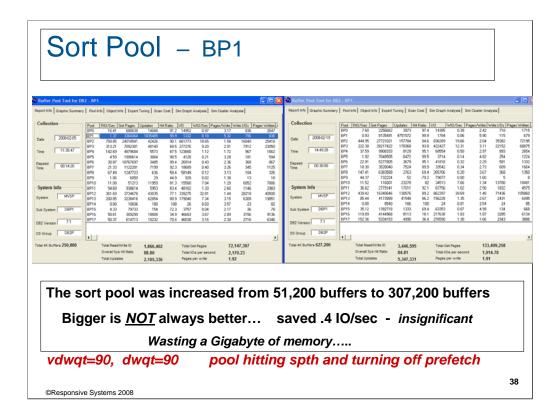
If we show the overall System Hit Ratio, it looks and sounds great. But the IO rate is quite high.... The is a lot of room for tuning to reduce the IO rate, save CPU, reduce transaction elapsed times, and improve productivity. The pools with the highest IO rates are quite a bit lower than that overall System Hit Ratio.



Effect of setting the sort pool thresholds <u>vdwqt and dwqt too high</u>. The pool hit spth, and prefetch was disabled.

Current Active Buffers – is a snapshot at the moment the Statistics record is produced, not a high water mark, and it's at 87% of the total pool buffers.

While performance looks good based on the Hit Ratios, the pool is having performance problems highlighted in Red. Other important metrics are highlighted in Blue.



Since the vdqwt and dwqt thresholds were incorrectly set at 90%, the pool was hitting spth, and turning off prefetch. An exception reporting system flagged the threshold problem, and suggested making the pool larger. It continued to hit spth for large sorts. The number of "buffers unavailable, or in-use" was generally 80% of the buffers, no matter how large the pool was made.

Sort Pool – Bigger is not always better

# Buffers	Read IO/Sec	Write IO/Sec	
51,200	1.37	.18	
256,000	1.25	.10	
307,200	0.93	.06	

The sort pool was oversized from the beginning, and memory would have been used more effectively on other pools. The reason for the expansive growth was....

Looking at one performance variable or exception may miss the real problem, and bad recommendations from an exception reporting product **You** need to understand performance, and evaluate recommendations...

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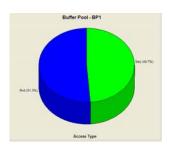
39

An analysis tool that only looked at the exception hitting spth, and merely recommended making the pool larger. No consideration was given to the vdwqt and dwqt settings, or the information that the pool consistently showed active (unavailable) pages at 80% of the pool.

Sort pool facts

- New pages are created in the pool, they are not read into the pool first — no initial read IO delay
- Pages that are written out, are not always read back
- Pools often have a high % of random getpages
- The prefetch read qty is 8 pages

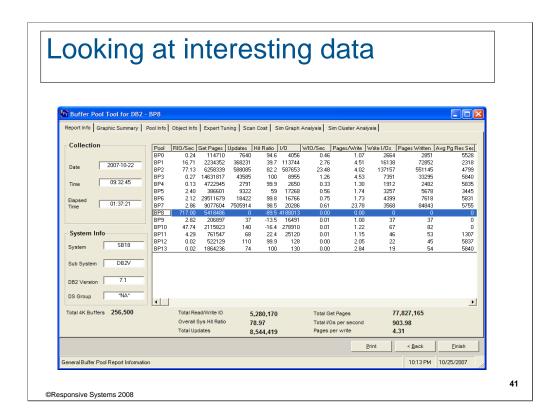
This is not an unusual access & usage illustration



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There is a lack of published information regarding sort functions and processing methods.

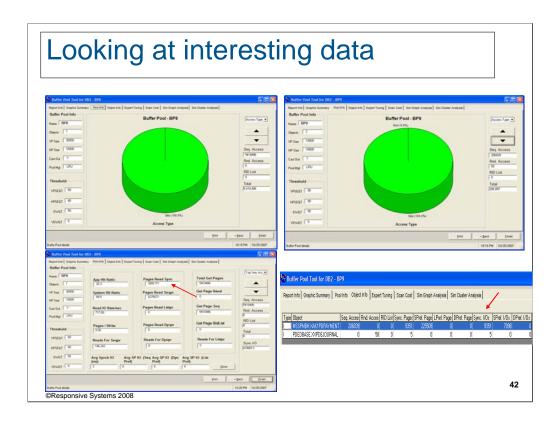


BP8 has the killer IO rate.

It has one object, a TS

BP9 has two objects, both indexes on the TS in BP8.

This is a heavy batch system, even though it is mid-morning.



BP8 has the killer IO rate.

It has one object, a TS

BP9 has two objects, both indexes on the TS in BP8.

Email to client after looking at data

You have a classic pool thrashing scenario, complicated by an application design/access issue.

Are you running multiple concurrent batch jobs against this, attempting to get better throughput by hitting different partitions?

BP8

As you said, one object, 10 partitions. The access to this object is 100% sequential. The high Synch IO tells me that the prefetched pages are thrown out *before a low priority batch job can be dispatched to read them*, and they are **re-read using synch IO**.

A major performance, and cpu/cost killer.

** you need to have the batch job(s) run at a higher system priority

** if you are running several, against different partitions, try running a few less concurrent jobs

BP9

Two indexes on the TS in BP8

PDPAYMENT is 100% SP access - for an index??? This is a design or sql coding problem, or lack of full stats for the index and TS..

It appears that there were at least 9,351 complete scans of this index.

The other index has almost no usage, and all random.

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Response from the client

Joel:

Dang you are good!! We had 3 jobs running going against 3 different partitions on the same table. I can't remember what SRVCLASS they were running in. Everybody has the ability to change SRVCLASS, which stinks, but that is my problem.

You say "9,351 complete scans of this index". I assume you are getting that number from this screen. How do you interpret this to be "complete scans" and not "pages read" since SYNC I/O's are one page at time?

Let's go back two slides and look at the data

This isn't genius, just experience.

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All the getpages issued were sequential.

So what causes a synch IO?

When a prefetch stream starts, it issues 1 synch IO for the first page, a SP for pages 2-32, and a SP for pages 33-64. When you hit page 32, it issues a SP for page 65-96, etc, etc

I am making an assumption here, since I am not looking at the actual trace data from the collection file, but since this is how sequential prefetch is initiated.... I'm assuming that each synch IO indicates the start of a scan...

Back to the client

All the getpages issued were sequential So - what causes a synch IO?

When a prefetch stream starts, it issues 1 synch IO for the first page, a SP for pages 2-32, and a SP for pages 33-64.

One synch IO, and two concurrent prefetch IOs.

When you hit page 32, it issues a SP for page 65-96, etc, etc

I am making an assumption here, since I am not looking at the actual trace data from the collection file, but since this is how sequential prefetch is initiated.... I'm assuming that each synch IO indicates the start of a scan...

I already covered this in the discussion of 3 slides back

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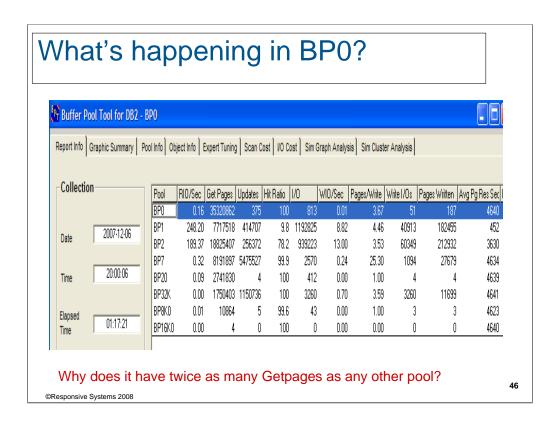
All the getpages issued were sequential.

So what causes a synch IO?

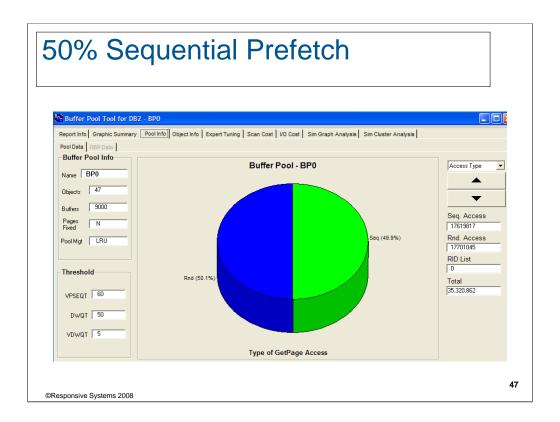
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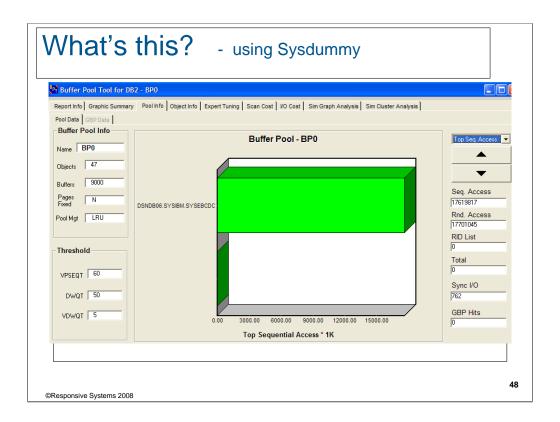
is initiated.... I'm assuming that each synch IO indicates the start of a scan...



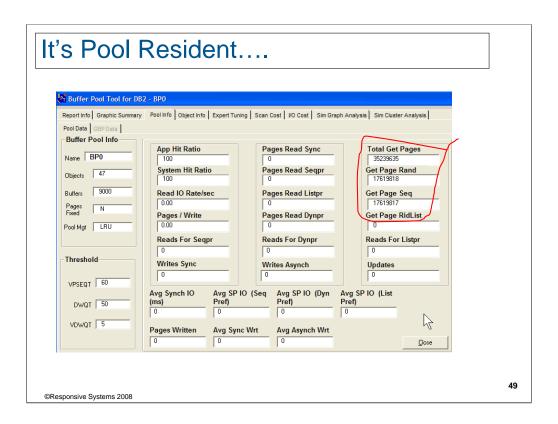
Very high number of Getpages, almost twice any other pool in the system. Application objects in the pool?



This is very unusual access for BP0.



Heavy application usage of SYSDUMMY



This object is "supposedly" one page in size. Updating catalog statistics for it did not get rid of the sequential activity.

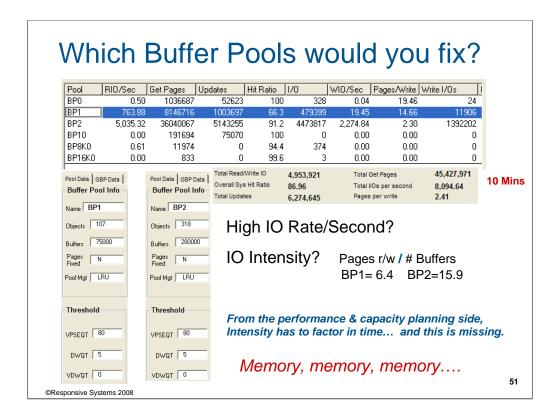
Pagefixing Buffer Pools

- Will save 8% of the IO CPU cost
- About a 15-20% reduction of DBM1 CPU cost
- Application savings are harder to measure, but not impossible
- You need REAL Memory availability before fixing memory. If the system starts to page, you die...

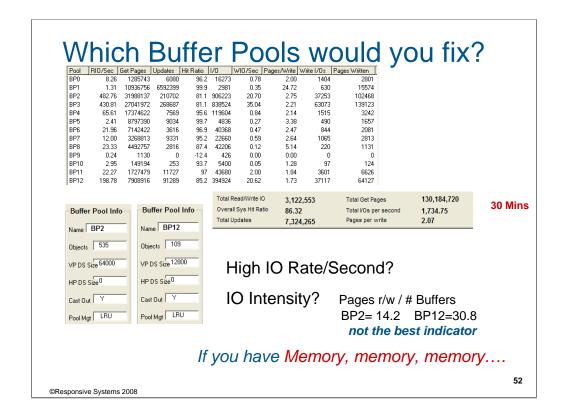
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Application saving would have to be measured from the 101 application accounting records, and accumulated over a period. Looking at a few records would not show anything.

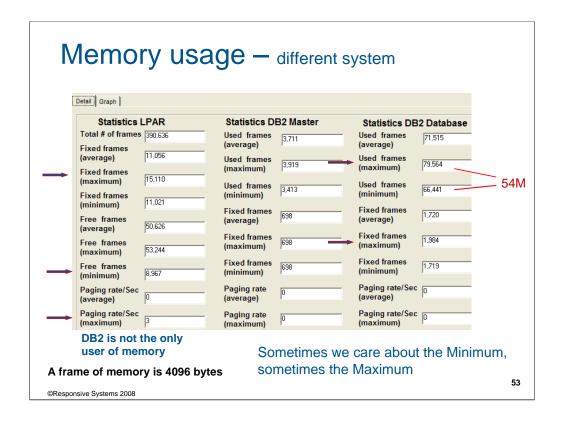


If you have the memory... the HIGH IO pool gives you the greatest saving, and CPU reduction.

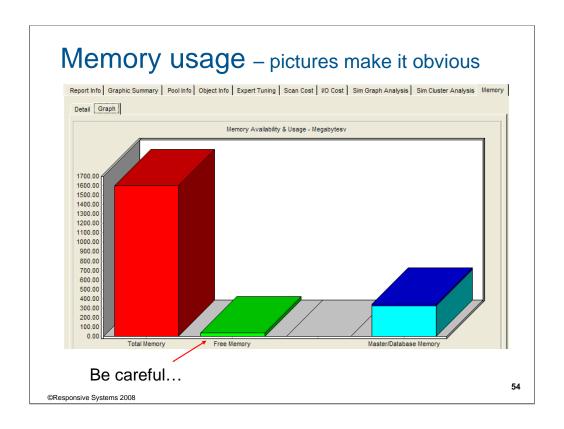


If you have the memory... the HIGH IO pool gives you the greatest saving, and CPU reduction. If you have the memory, BP2 will give you much better payback than BP12.

So, IO Intensity is **NOT** the best indicator.



If you have the memory... the HIGH IO pool gives you the greatest saving, and CPU reduction. But you don't dare pagefix anything when available memory is this low, and yu already see a count for system paging.



If you have the memory... the HIGH IO pool gives you the greatest saving, and CPU reduction. But you don't dare pagefix anything when available memory is this low.

Tracking & comparing performance....

 One of the largest issues we face, is trying to determine if performance and workloads are similar or

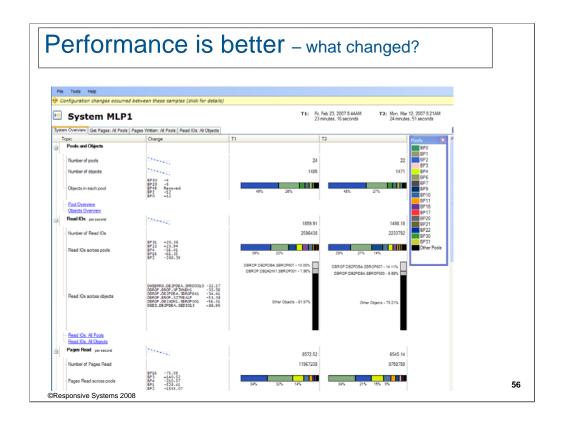
We know there has been a change in performance, so what changed?

- What and where is the difference?
 - Application workload?
 - System changes?
 - Both?

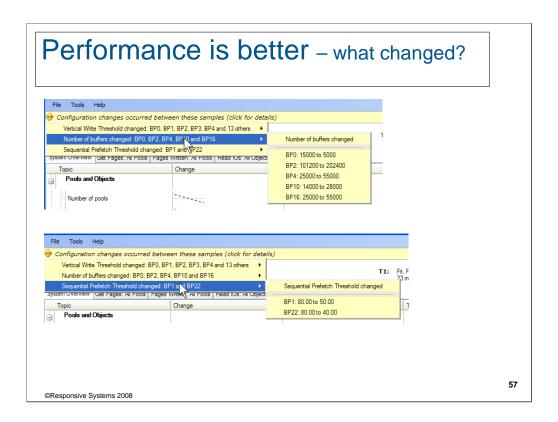
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Some performance differences are obvious, but the reasons may be difficult to find. In most cases, we are looking at performance degradations, and need to find out why the users are complaining about poor response times, or batch jobs are running too long. An automated way of looking at two sets of data, and comparing performance, and highlighting the differences – but at the system level, and at the workload/object usage level would make life much easier.



We have both configuration and workload changes. Objects not accessed in the second run, a pool not utilized, pool thresholds have been changed.



This highlights the system configuration changes from the earlier run.

DB2 V8 and V9

- Exploitation of 64bit memory, and moving more control blocks out of the DBM1 address space
- Access to many gigabytes of memory for pools
 - · Does not change basic pool tuning methodologies
 - Grouping by Random and Sequential, and then by working set size is the industry proven technique (RAMOS, SAMOS)
 - Working set size has no relationship to number of pages for an object (catalog statistics)
 - An object has 1,000,000 pages, but the maximum number in the pool during a specific time period is 2,967
 - The wkset is 2,967
 - · This may change, as the pool size is increased or decreased

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The working set size of an object is the number of pages in the pool at a given point in time. There is no relationship between the working set size, and the information you will find in the DB2 catalog.

- More areas of memory moved above the 2Gig Bar
 - Large parts of the EDM Pool
 - SKCT and SKPT
 - DBDs
 - · Parts of CT and PT
 - Parts of Dynamic SQL

You need REAL memory behind everything...

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A very large memory relief for big systems. Keep in mind that you still need the REAL memory available on the machine!!

- Increasing prefetch quantities and using larger page sizes to reduce IO (vpseqt * #Buffers) > 40,000
 - Double the prefetch and deferred write number of buffers
 - IO remains a huge throughput concern for system scalability
- Supports buffer pool sizes > 5 Gigabytes
 - Up to 1 Terabyte of memory for DB2
- Increased usage of 32K sort file to reduce IO
 - · Larger and multiple 32K sort objects
 - Improved sort logic row < 101 bytes uses 4K longer uses 32K
 - It's possible to limit the amount of sort space for a thread (zparm)

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IO remains a primary concern for system scalability, and these concerns are being addressed in every new version of DB2.

- Automatic Pool Size Management option autosize=yes
 - Function integrated with WLM
 - Can increase/decrease pool size by 25% of initial size
 - Increments? Not much real information available yet
 - Based on long term trends
 - This is not defined anyplace
 - Long term performance data is a major problem...
 - Tries to take the memory from other low activity pools first
 - Based on a random hit ratio hit ratios are not valid as a performance metric
 - Seeming fallacy of this approach lack of <u>IO</u> prediction capability
 - We already proved bigger is not always better
 - Will WLM reduce it if the increase does not improve performance?

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There are good long term possibilities for this type of approach, and it's obvious that this is just a first cut implementation. It remains to be seen if this provides any real benefit. I suspect it may for small to medium systems at installations where there isn't a lot of DB2 performance knowledge. Large and high performance systems still need real tuning expertise.

Also, remember that the effective way to get good performance is through the proper grouping of objects (Ramos/Samos), and not by throwing memory at few large pools.

- What changes in regard to pool tuning?
 - Absolutely nothing
- Methodologies remain the same Ramos/Samos & Wkset sizes
- There are opportunities for ever larger pools, if you have the memory.... and are sure you get a real benefit
 - Remember, paging is to DASD, not expanded memory, so that's 1,000 times slower
 - Just as previous environments, if you start to page, your pool
 performance may look better statistically, but the users performance
 is worse... be really sure memory is available....
 - WLM is supposed to manage memory better, and it can, if it's set up properly

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The basics of performance tuning have not changed over the last four decades, and certainly won't change over the next decade. CPU, Memory, and IO are the important tuning metrics.

DB2 Version 9 – some Important Zparms

- SJMXPOOL Memory pool for Starjoins, must be enabled by the STARJOIN zparm. Up to 1 Gig of memory, above the 2 Gig bar
- MXTDCACH in memory data caching, other than Starjoin, allocated per thread
- DSVCI allows DB2 to create datasets with a CI size that matches the page size, such as 8K, 16K, 32K
- MAXTEMPS the max workfile storage an agent can use

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While there are a great many zparm changes in V9, and they are ALL important to your system, these are a few you may want to start with.

DB2 Version 9 – overhead comparison

Table 4-2 Tabulated data for relative CPU % increase for 100 package query application

Accounting Class	CPU%			
	DB2 V8 CL1	DB2 V8 CL2	DB2 V9 CL1	DB2 V9 CL2
1	-	-	-	-
1,2	2.87	-	3.52	-
1,2,3	3.18	0.15	3.96	0.28
1,2,3,7,8	6.96	4.23	5.91	2.43
1,2,3,7,8,10	23.29	21.98	19.82	17.30

The tabulated data shown in Table 4-2 is also shown graphically in Figure 4-13.

Trace filtering capability provides an overhead reduction

From DB2 9 for z/OS Performance Topics Redbook

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So – Red is bad, Green is good. The cost of the base C1, C2, C3 accounting traces has gone up. The cost of adding package level information has decreased a lot.

Classes 7, 8, 10 are the high overhead classes, just as V8, but substantially reduced. The traces can have $\underline{Include}$ and $\underline{Exclude}$ lists by many criteria, such as

USERID, WRKSTN, APPNAME, PKGLOC, PKGCOL, PKGPROG, CONNID, CORRID, and ROLE.



Systems and tools are becoming better all the time. While there are some attempts at self tuning systems, or some parts that have a potential for self tuning, we're a long way from realizing these goals. The largest obstacle is the overhead cost of effective approaches. Doing it properly is a lot of work, and very expensive from a CPU perspective.

So – it still requires some work to achieve better performance, and the rewards of lower operational costs that better performance provides.



Session A03

DB2 System Performance, the Basics... and a Bit, to a Lot More

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IDUG 2008 North America